

**INTERNATIONAL CENTRE FOR APPLIED SCIENCES**  
 (Manipal University)  
**III SEMESTER B.S. DEGREE EXAMINATION –MAY 2016**  
**SUBJECT: FORMAL LANGUAGES AND AUTOMATA THEORY (CS 233)**  
 (BRANCH: COMPUTER SCIENCE)  
**24<sup>TH</sup> MAY, 2016**

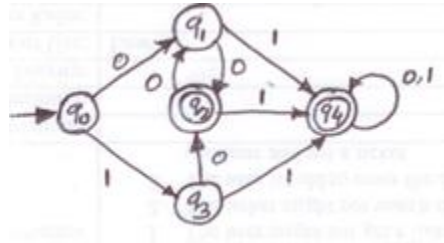
Time: 3 Hours

Max. Marks: 100

- ✓ **Answer ANY FIVE full Questions.**
- ✓ **Missing data, if any, may be suitably assumed.**

- 1A.** Define finite automata. Explain deterministic and non-deterministic finite automata. with the help of an example for each . **6**
- 1B.** Design a DFA which contains exactly one 'a' and exactly two b's. **4**
- 1C.** Prove that for any transition function  $\delta$  and for any two input strings x and y,  

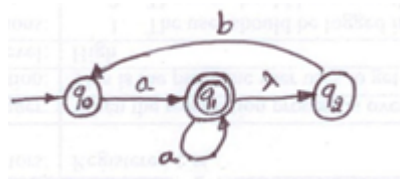
$$\delta^*(q, xy) = \delta(\delta^*(q, x), y)$$
 **6**
- 1D.** Obtain grammar for regular expression representing strings of a's and b's, having a substring 'ab'. What language does this grammar represent? **4**
- 2A.** Minimize the following DFA Fig-2A. Clearly show all the steps.



**(Fig-2A)**

**8**

- 2B.** Convert following NFA to DFA Fig-2. Show all the steps.



**(Fig-2B)**

**6**

- 2C.** Obtain regular expression for the following sets
- $L_1 = \{ a^n b^m : (n+m) \text{ is even} \}$
  - $L_2 = \{ \text{Strings of a's and 'b's ending with 'ab' or 'ba'} \}$
- 6**
- 3A.** State pumping lemma for regular languages with rules and use it to check whether the language  

$$L = \{ 1^p \mid \text{where } p \text{ is prime} \}$$
 is regular or not. **8**

- 3B.** Define right linear and left linear grammar with example for each .Construct a left linear grammar by designing respective non deterministic automata for the given languages. 12
- (i)  $L \{(aab)^*ab\}$  (ii)  $L(abb^*)$
- 4A.** Find context free grammar for the language  $L = \{ a^n b^m c^k : m = n \text{ or } m \leq k \}$  4
- 4B.** Given the set of production  
 $S \rightarrow S^*A \mid A \quad A \rightarrow A+B \mid B \quad B \rightarrow (S) \mid a \mid b$   
 Find the left most derivation and draw the corresponding derivation tree for the string  $a^*(b+a)$ . 6
- 4C.** Find S grammar for the language  $L = \{ a^n b^{n+1} \mid n \geq 2 \}$  5
- 4D.** Test whether the following grammar is ambiguous or not  
 $S \rightarrow AB \mid aaB \quad A \rightarrow aA \mid a \quad B \rightarrow b$   
 If it is ambiguous construct an equivalent unambiguous grammar. 5
- 5A.** Eliminate useless,  $\lambda$  and unit productions from the following  
 $S \rightarrow abAB \mid A \quad A \rightarrow bAB \mid \lambda \quad B \rightarrow BAa \mid A \mid \lambda$  6
- 5B.** Convert the following into CNF:  
 $S \rightarrow ASB \mid \lambda \quad A \rightarrow aAS \mid a \quad B \rightarrow SbS \mid A \mid bb$  6
- 5C.** Convert the following into GNF:  
 $S \rightarrow AA \mid 0 \quad A \rightarrow SS \mid 1$  4
- 5D.** Show the following two grammar are equivalent  
 (i)  $S \rightarrow abAB \mid ba$  (ii)  $S \rightarrow abAaA \mid abAbb \mid ba$   
 $A \rightarrow aaa$   $A \rightarrow aaa$   
 $B \rightarrow aA \mid bb$  4
- 6A.** With neat figure explain the language families in the Chomsky Hierarchy 8
- 6B.** Construct NPDA for the language  $L = \{ ww^R \mid w \in \{a,b\}^+ \}$  6
- 6C.** Convert the following CFG to PDA  
 $S \rightarrow 0A$   
 $A \rightarrow 0AB \mid 1$   
 $B \rightarrow 1$  6
- 7A.** Design nondeterministic pushdown automata with its transition graph for the following grammar. Trace it for the string  $w = aaabc$ . 10
- $S \rightarrow aA$   
 $A \rightarrow aABC \mid bB \mid a$   
 $B \rightarrow b$   
 $C \rightarrow c$
- 7B.** Design Turing machine which accepts the language  $L = \{ 1^n 2^n \mid n \geq 1 \}$ . Draw transition graph and also trace it for  $w = 1122$ . 10
- 8A.** Define recursively enumerable and recursive languages .Is the family of recursive languages closed under concatenation? 6
- 8B.** Discuss the concept of context sensitive grammar and language and also give one example for the same. 4
- 8C.** Give formal definition of Off- line turing machine with the help of figure. 4
- 8D.** Suppose we make a restriction that a turing machine must always write different from the one it reads, that is if  $\delta(q_i, a) = (q_j, b, L \text{ or } R)$ , then a and b must be different. Does this limitation reduce the power of automaton? Justify. 4

