

- d) $g(x, h(y, z), d)$
e) $f(f(g(d, x), f(g(d, x), y, g(y, d)), g(d, d)), g(f(d, d, x), d), z)$.
- ii) The length of a term over \mathcal{F} is the length of its string representation, where we count all commas and parentheses. List all variable-free terms over \mathcal{F} of length less than 10.
- iii) The height of a term over \mathcal{F} is defined as 1 plus the length of the longest path in its parse tree. List all variable-free terms over \mathcal{F} of height less than 4.
- 2C. Draw the parse tree for the following CTL formulas:
- i) $AG(q \rightarrow EGr)$
ii) $A[p U EFr]$
iii) $EFEGp \rightarrow AFr$
iv) $AG(p \rightarrow A[p U (\neg p \wedge A[\neg p U q]))]$.

[5+3+2]

3A. Let Γ be the set of formulas

$$\Gamma = \left\{ C(p_1 \vee p_2 \vee p_3), \right. \\
C(p_1 \rightarrow K_2 p_1), C(\neg p_1 \rightarrow K_2 \neg p_1), \\
C(p_1 \rightarrow K_3 p_1), C(\neg p_1 \rightarrow K_3 \neg p_1), \\
C(p_2 \rightarrow K_1 p_2), C(\neg p_2 \rightarrow K_1 \neg p_2), \\
C(p_2 \rightarrow K_3 p_2), C(\neg p_2 \rightarrow K_3 \neg p_2) \\
C(p_3 \rightarrow K_1 p_3), C(\neg p_3 \rightarrow K_1 \neg p_3) \\
\left. C(p_3 \rightarrow K_2 p_3), C(\neg p_3 \rightarrow K_2 \neg p_3) \right\}.$$

Prove the sequent $\Gamma, C(p_2 \vee p_3), C(\neg K_2 p_2 \wedge \neg K_2 \neg p_2) \vdash K_3 p_3$.

- 3B. Consider a formula scheme $\Box\phi \rightarrow \Box\Box\phi$, as per the correspondence theory prove that its accessibility relation R has transitive property.
- 3C. Consider the Kripke model \mathcal{M} depicted in Figure Q.3C. For each of the following determine whether it holds:

- i) $a \Vdash q$
ii) $a \Vdash \Box\Box q$
iii) $a \Vdash \Box\Diamond\neg q$
iv) $c \Vdash \Box\perp$.

[5+3+2]

4A. Consider the wise-men puzzle. Justify your answers.

- i) Each man is asked the question 'Do you know the color of your hat?' Suppose that the first man says 'no,' but the second one says 'yes.' Given this information together with the common knowledge, can we infer the color of his hat?
- ii) Can we predict whether the third man will now answer 'yes' or 'no'?

iii) What would be the situation if the third man were blind?

4B. Prove sequent $\vdash_K \Box(p \rightarrow q) \vdash \Diamond p \rightarrow \Diamond q$.

4C. Using the natural deduction rules for $KT45^n$, prove the validity of $K_i Cp \leftrightarrow Cp$.

[5+3+2]

5A. Give Backus Naur form for branching tree logic and define the relation $\mathcal{M}, s \models \phi$ by structural induction on ϕ .

5B. Consider the system of Figure Q.5B. For each of the formula ϕ :

- i) Ga
- ii) $a U b$
- iii) $a U X(a \wedge \neg b)$
- iv) $X\neg b \wedge G(\neg a \vee \neg b)$
- v) $X(a \wedge b) \wedge F(\neg a \wedge \neg b)$.

Find a path from the initial state q_3 which satisfies ϕ .

5C. Prove the sequent $\forall x \forall y P(x, y) \vdash \forall u \forall v P(u, v)$.

[5+3+2]

6A. Does $\models \phi$ hold for the ϕ below? Justify your answer.

- i) $(p \rightarrow q) \vee (q \rightarrow r)$
- ii) $((q \rightarrow (p \vee (q \rightarrow p))) \vee \neg(p \rightarrow q)) \rightarrow p$.

6B. Describe basic labelling algorithms with reference to CTL model checking.

6C. You are asked to prove the validity of a sequent $\forall x(P(x) \rightarrow Q(x)), \exists x P(x) \vdash \exists x Q(x)$. Suppose your proof is as follows:

1.	$\forall x(P(x) \rightarrow Q(x))$	Premise
2.	$\exists x P(x)$	Premise
3.	$x_0, P(x_0)$	Assumption
4.	$P(x_0) \rightarrow Q(x_0)$	$\forall x e, 1$
5.	$Q(x_0)$	$\rightarrow e, 3, 4$
6.	$Q(x_0)$	$\exists x e, 2, 3-5$
7.	$\exists x Q(x)$	$\exists x i, 6$

Is this proof correct? If not then what is wrong with this proof?

[5+3+2]

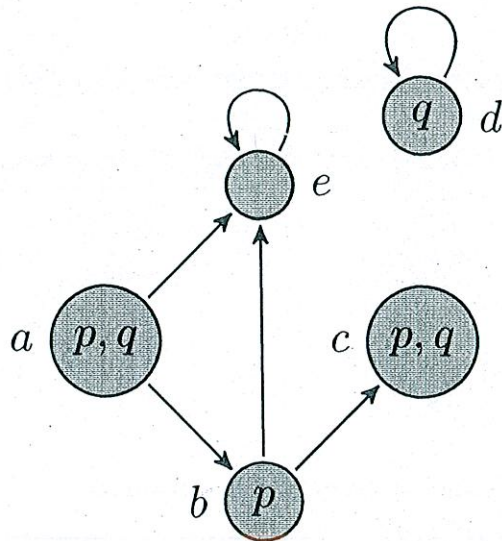


Figure: Q.3C

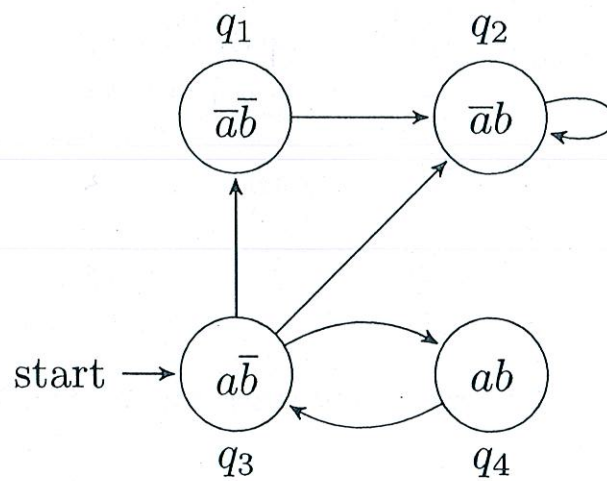


Figure: Q.5B